## Codes with Prescribed Permutation Group

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In this paper a program is proposed how to determine codes with given transitive permutation group. The approach is module theoretic, based on a study of monomial actions and projective representations. Some highly transitive groups are discussed in detail.

There appear slightly different concepts of (linear) codes in the literature. Following Ward [16] and Rasala [11] a code over some commutative ring  $\overline{F}$ with unity will be a triple  $(V, B, C)$ , where V is a free F-module of finite rank with basis  $R$  and submodule  $C$ . By convention we then call  $C$  a code having ambient space V and ambient basis B. F is the alphabet of C, the rank  $n$  of V its length, and C is an  $(n, k)$ -code if C is free of rank k. (In this paper F will always be a field.)

The Hamming weight of a vector (word) in  $V$  is the cardinality of its support with respect to the given basis. The minimum weight of a code  $C$  is a measure for its error-correcting capability. Hence morphisms between codes should preserve the Hamming weight. This leads to the definition: A morphism  $(V, B, C) \rightarrow (V', B', C')$  of codes over F is an injective F-linear map  $\mu: V \to V'$  with  $C\mu \subseteq C'$  sending any  $e \in B$  to a scalar multiple of some  $e' \in B'$ . The codes C and C' are isomorphic if  $\mu$  is bijective and  $C\mu = C'$ .

 $ML(C)$  denotes the group of all (code) automorphisms from  $(V, B, C)$ onto itself, the *monomial linear group* of C.  $(ML(C))$  can be represented, with respect to B, by monomial matrices.) Let  $B = (e_i)$ ,  $i \in \Omega = \{0, ..., n-1\}$ . Every  $\mu \in ML(C)$  determines a permutation  $\bar{\mu}$  on  $\Omega$  by  $e_i \mu \in \langle e_{i\bar{\mu}} \rangle$ . The map  $\mu \mapsto \bar{\mu}$  is an epimorphism of  $ML(C)$  onto a subgroup  $PML(C)$  of the symmetric group on  $\Omega$ . PML(C) is called the *permutation group* of C. C admits a permutation group G on  $\Omega$  if G is a subgroup of PML(C). The elements of ker( $u \mapsto \bar{u}$ ) are the diagonal automorphisms of C.

Observe that the transitivity behaviour of the permutation group of a code is a measure for its homogeneity. Codes having (multiply) transitive permutation groups have good error-correcting properties and provide for powerful decoding methods. Actually there are many interesting codes with a

multiply transitive permutation group. For instance the extended quadratic residue (*QR*-) code of length  $p + 1$ , p an odd prime, admits PSL(2, p). (A generalization to symplectic groups can be found in Ward [16].) The 5transitive Mathieu groups  $\mathfrak{M}_{12}$  and  $\mathfrak{M}_{24}$  are the permutation groups of the extended ternary and binary Golay codes, respectively. The binary Reed-Muller codes of length  $2<sup>m</sup>$  admit the affine groups Aff $(2<sup>m</sup>, 2)$ . (A general reference for this is MacWilliams and Sloane [9].)

For theoretical and practical reasons one may ask for a method to determine all codes with prescribed permutation group. In attacking this problem we show, via some kind of Krull-Schmidt decomposition for codes, that it suffices to construct those codes whose diagonal automorphisms are scalar multiplications (see Section 1). Every diagonal automorphism of a nontrivial code C is scalar, for instance, if  $PML(C)$  is primitive (Theorem 1.3). In this case we have a central group extension

$$
F^* \rightarrow ML(C) \rightarrow PML(C),
$$

where  $F^*$  denotes the multiplicative group of the field F.

A group E is said to act (monomially) on a code  $(V, B, C)$  if there is given a homomorphism  $E \to ML(C)$ . E induces a permutation group G on  $\Omega$ , the index set of B. In the situation  $(*)$  we obtain a projective representation of G on V which lifts back to the given (ordinary) representation of  $E$  on V. Under suitable assumptions, this projective representation can be lifted also by stem covers of  $G$  ("Darstellungsgruppen" in Schur's terminology). Moreover, if G acts transitively on  $\Omega$  and  $E_0$  is the subgroup of E fixing  $U = \langle e_0 \rangle$ , then the monomial action of E on V can be replaced by that induced on  $U^E = U \otimes_{E_0} FE$  (Proposition 2.1).

This will serve as a principle for constructing codes admitting a given primitive permutation group  $(G, \Omega)$ : Suppose E is a stem cover of G and  $E_0$ is the inverse image in  $E$  of a point stabilizer. Then the  $FE$ -submodules of all induced modules  $U^E$ , U being a 1-dimensional  $FE_0$ -module, yield a complete list of codes over F admitting  $(G, \Omega)$ , provided  $Ext(G/G', F^*) = 0$ (Theorem 3.1). This condition is fulfilled, for instance, if  $F$  is algebraically closed or  $G = G'$  is perfect. In general one can start with an algebraically closed field of scalars, which is appropriate also for module theoretic reasons. Then one has to find, for any submodule C of  $U^E$ , the smallest fields of realization.

To illustrate the program we will determine all codes admitting alternating groups or Mathieu groups. It turns out that the alternating groups  $\mathfrak{A}_n$  of degree  $n \geq 7$  occur only in the permutation group of the repetition code and its dual (Theorem 4.4). Here we make use of Schur's work  $[12]$  on the multipliers of alternating groups. The Mathieu groups only leave invariant Golay codes, besides the repetition code and its dual. This depends on results of Burgoyne and Fong [2] (and P. Mazet [18]) on the Schur multipliers of the Mathieu groups.

The paper is concluded by a discussion of  $QR$ -codes. We show that the extended QR-codes (of length  $p + 1$ ) are characterized by the property that they admit  $PSL(2, p)$  but not  $PGL(2, p)$  (Theorem 6.2). It is conjectured that the (full) permutation group G of such a code is precisely  $PSL(2, p)$ provided  $p > 23$ . We can prove, at least, that G is a proper subgroup of  $\mathfrak{A}_{p+1}$ if  $p > 5$ . This answers a conjecture by Rasala [11] to the affirmative. If  $p > 23$  and  $G \neq PSL(2, p)$ , then G would be an "unknown" simple group being 4-transitive on  $p + 1$  letters (Theorem 6.4).

### 1. INDECOMPOSABLE CODES

Let C be a code over F with ambient basis  $B = (e_i)$ ,  $i \in \Omega = \{0, ..., n-1\}$ . If  $B' \subseteq B$  then  $C' = C \cap \langle B' \rangle$  is regarded as a code with ambient space  $\langle B' \rangle$ and ambient basis  $B'$ . C is called decomposable if B can be partitioned into at least two nonempty subsets  $B_j$  such that  $C = \bigoplus C_j$ , where  $C_j = C \cap \langle B_j \rangle$ , and indecomposable otherwise. There is a unique partition of  $\vec{B}$  into subsets  $B_j$  such that  $C = \bigoplus C_j$  and each  $C_j$  is indecomposable (Krull–Schmidt).

The decomposition of C into its indecomposable components  $C_i$  can be studied from a different point of view. Call a nonzero vector  $v \in C$ indecomposable if v is not the sum of two nonzero vectors in C with disjoint supports. Every vector is a sum of indecomposables which, however, are not uniquely determined. If  $d$  is the minimum weight of  $C$ , then any nonzero vector in C of weight at most  $2d - 1$  is indecomposable. (Recall that the weight of  $v = \sum a_i e_i$  is the cardinality of supp $(v) = \{i | i \in \Omega, a_i \neq 0\}$ .

Define the binary relation  $A = A_C$  on  $\Omega$  to be the set of all pairs  $(i, j) \in \Omega^2$ such that there is an indecomposable  $v \in C$  having i, j in its support. Let  $\overline{A} = \overline{A}_C$  be the smallest equivalence relation on  $\Omega$  containing A. Then  $(i, j) \in \overline{A}$  if and only if  $i = j$  or there are indecomposable  $v_k \in C$   $(1 \leq k \leq m)$ such that  $i \in \text{supp}(v_1)$ ,  $j \in \text{supp}(v_m)$ , and  $\text{supp}(v_{k-1}) \cap \text{supp}(v_k) \neq \emptyset$  for  $2 \leq k \leq m$ . Note that  $A_c$  and  $\overline{A}_c$  are invariant under the automorphism group  $ML(C)$ , i.e., under  $PML(C)$ .

(1.1) LEMMA. Let  $(B_j)$  be the partition of B corresponding to the equivalence classes of  $\overline{A}_C$  and  $C_j = C \cap \langle B_j \rangle$ . Then  $C = \bigoplus C_j$  is the decomposition of  $C$  into its indecomposable components.

*Proof.* If  $v \in C$  is indecomposable, supp(v) is contained in just one equivalence class of  $\overline{A}$ . Since any  $v \in C$  is a sum of indecomposable vectors in  $C$ , it is enough to show that each  $C_j$  is indecomposable. Assume  $J_j = B' \cup B''$  (disjoint,  $B' \neq \emptyset \neq B''$ ) and  $C_j = (C_j \cap$ 

Since  $B_i$  corresponds to an equivalence class of  $\overline{A}$ , there must be an indecomposable vector  $v \in C_i$  such that supp(v) meets the index sets of both B' and B", which is impossible.  $\blacksquare$ 

The indecomposable components of  $C$  can be related to the structure of ML(C). To explain this we introduce a further equivalence relation  $\Delta = A_c$ on  $\Omega$ . Let  $(i, j) \in \Delta$  if each diagonal automorphism of C multiplies both  $e_i$ and  $e_i$  with the same scalar. Of course, if  $F = \mathbb{F}_2$  then  $\Delta$  is the universal relation on  $\Omega$ . Thus  $\Delta$  is interesting only when  $F \neq \mathbb{F}_2$ .

(1.2) THEOREM. Suppose  $F \neq \mathbb{F}_2$ . Then  $\overline{A}_C$  and  $A_C$  coincide. In particular, C is indecomposable if and only if every diagonal automorphism of C is a scalar multiplication.

*Proof.* From Lemma 1.1 it follows  $\Delta \subseteq \overline{A}$ , because of  $|F| > 2$ . To prove the converse let  $(B_i)$  be the partition of B associated to the equivalence classes of  $\Delta$ . Let  $v \in C$ . Then there are unique  $v_i \in \langle B_i' \rangle$  such that  $v = \sum v_i$ . We claim that all  $v_j$  belong to C. Define  $m_v = \max\{j \mid v_j \neq 0\}$ ,  $m_0 = 0$ . The claim is obvious if  $m_v \le 1$ . Let  $m_v = m > 1$ . Fix *j* between 1 and  $m - 1$ . By definition of  $\Delta$  there exists a diagonal automorphism  $x$  of  $C$  such that  $v_k x = a_k v_k$   $(1 \leq k \leq m)$  and  $a_i \neq a_m$ . Then

$$
w = a_m v - vx = \sum_{k=1}^{m-1} (a_m - a_k) v_k
$$

is in C and satisfies  $m_w \le m - 1$ . By induction  $w_j = (a_m - a_j) v_j \in C$ , hence  $v_j \in C$ . Also  $v_m = v - \sum_{k=1}^{m-1} v_k \in C$ , as claimed.

We have established that, for every indecomposable  $v \in C$ , supp $(v)$  is completely contained in some equivalence class of  $\Delta$ . Therefore  $\Lambda \subseteq \Delta$ , hence also  $\bar{\Lambda} \subseteq \Lambda$ . I

We now give a sufficient condition for a code to be indecomposable in terms of its permutation group.

(1.3) THEOREM. If C is a nontrivial code such that  $PML(C)$  is primitive, then C is indecomposable and every diagonal automorphism of C is scalar.

*Proof.* The minimum weight  $d$  of  $C$  is at least 2, by transitivity of  $G = PML(C)$ . Consequently  $\overline{A}$  is not the diagonal in  $\Omega^2$ . Since  $\overline{A}$  is Ginvariant and G is primitive, it follows that  $\overline{A}$  is the universal relation on  $\Omega$ . By  $(1.1)$  C is indecomposable. Finally apply Theorem 1.2.

Observe that there are decomposable codes having a transitive permutation group, e.g.,  $V = F^4$ ,  $C = \langle 1010 \rangle \oplus \langle 0101 \rangle$ . Here C is cyclic. However, we have the following criterion.

 $(1.4)$  PROPOSITION. Let C be cyclic. If C contains an indecomposable vector v such that there are i, j in supp(v) with  $i - j$  coprime to the length n of  $C$ , then  $C$  is indecomposable.

## *Proof.* Straighforward.  $\blacksquare$

The structure of a code C is completely determined by the structure of its indecomposable components. So, in principle, we may restrict our attention to the study of indecomposable codes.

### 2. ACTIONS ON CODES AND INDUCED MODULES

Let  $(V, B, C)$  be an  $(n, k)$ -code over F,  $B = (e_i), i \in \Omega = \{0, ..., n-1\}.$ Suppose we have a group homomorphism  $\varphi: E \to ML(C)$ . Then E is said to act (monomially) on C (via  $\varphi$ ). Composing  $\varphi$  and the natural epimorphism from  $ML(C)$  onto PML(C) yields a map  $E \to PML(C)$  whose image G is a permutation group on  $\Omega$ .

Clearly V is an FE-module via  $\varphi$ , with invariant subspace C. Assume E (i.e., G) is transitive on  $\Omega$ . Let  $E_0$  be the largest subgroup of E leaving invariant the 1-dimensional subspace  $U = \langle e_0 \rangle$ . Then, for each  $i \in \Omega$ , there exists  $x_i \in E$  mapping U onto  $\langle e_i \rangle$ . Hence

$$
V = \bigoplus U x_i
$$

is an FE-module induced by the  $FE_0$ -module U.

(2.1) PROPOSITION. Assume E acts on  $(V, B, C)$  and is transitive on  $\Omega$ . Let  $E_0$  be the subgroup of E leaving invariant  $U = \langle e_0 \rangle$  and let  $V' = U \otimes_{E_0} FE$ . Choose a right transversal  $(x_i)$  to  $E_0$  in E indexed such that  $e_0x_i \in \langle e_i \rangle$ , say  $e_0x_i = a_ie_i$ . Let  $e'_i = e_0 \otimes x_i$  and  $B' = (e'_i)$ . Then the linear map  $\mu: V \to V$  given by  $e_i' \mapsto a_i e_i$  is a monomial isomorphism  $(V', B') \rightarrow (V, B)$  of FE-modules, and the preimage C' of C represents a code  $isomorphic$  to  $C$ .

*Proof.* It is immediate that  $B' = (e'_i)$  is a basis for V'. E operates on V' by

$$
e'_ix = (e_0 \otimes x_i)x = c_ie_0 \otimes x_j = c_ie'_j,
$$

where  $x_i x = \bar{x}_i x_i$ , with  $\bar{x}_i \in E_0$  and  $e_0 \bar{x}_i = c_i e_0$ . Since also

$$
(a_i e_i) x = (e_o x_i) x = c_i e_o x_j = c_i (a_j e_j),
$$

 $\mu$  is an FE-isomorphism. We are done. 體

Identifying  $V$  and  $V'$  in the situation of Lemma 2.1 is now justified. The code C can be represented as a submodule of the induced module  $U^{\varepsilon} = U \otimes_{E_0} F E$ ,  $U = \langle e_0 \rangle$ , equipped with a basis  $B = (e_i)$ , where  $e_i = e_o \otimes x_i$ for some right transversal  $(x_i)$  to  $E_0$  in E, choosing  $x_0 = 1$ . This notation is fixed in the sequel.

 $(2.2)$  Remark. In theory, we may construct all codes over F admitting a transitive permutation group  $(G, \Omega)$  as follows: Consider any group extension  $A \rightarrow E \rightarrow G$  with A abelian. Let  $E_0$  be the inverse image in E of a point stabilizer  $G_0$ . Inducing up to E all 1-dimensional  $FE_0$ -modules U we obtain all (transitive) monomial representations of  $E$  with permutation group  $(G, \Omega)$ . The codes admitting  $(G, \Omega)$  occur as submodules of all  $V = U^F$ .

By a monomial action of an extension  $E$  of a (transitive) permutation group  $(G, \Omega)$  we always mean a monomial representation of E inducing  $(G, \Omega)$ .

In general, it is fairly hopeless to construct all required group extensions of G. However, in order to obtain all those codes which are indecomposable we have to consider only the case where A is central in E and isomorphic to a subgroup of  $F^*$  (Theorem 1.2). Moreover, any supplement to A in E will leave invariant the same subspaces. When  $F$  is finite we may take minimal supplements yielding central Frattini extensions of G. There exists, to any finite group G and any finite field  $F$ , a unique maximal (central) Frattini extension  $A_F \rightarrowtail G_F \rightarrowtail G$ , with  $A_F$  of exponent dividing  $|F^*|$ , having any other such extension of G as epimorphic image over  $G$  (i.e., inducing the identity on  $G$ ). This is a slight generalization of a classical result by Gaschütz  $[5]$ . We omit the details. In Section 3 we will see that one can use without loss stem covers of G instead.

We present some basic facts concerning monomial actions and induced modules. Throughout let  $E$  be a finite extension of the transitive permutation group  $(G, \Omega)$  and  $V = U^E$  for some 1-dimensional  $FE_0$ -module U, the basis  $B = (e_i)$  of V indexed by  $\Omega$ . (If W is an FH-module and  $\alpha$  is (or induces) an automorphism of H, then  $W_a$  is the vector space W with module structure  $w \circ x = wx^{\alpha}$  for  $w \in W$ ,  $x \in H$ .)

(2.3) LEMMA. Assume G acts 2-transitively on  $\Omega$ . Then the F-dimension of  $\text{End}_{\mathcal{F}}(V)$  is at most 2, and it is 1 if and only if the restriction W of U to  $H = E_0 \cap E_0^{y^{-1}}$ , for any  $y \in E - E_0$ , is not isomorphic to  $W_y$ .

*Proof.* Apply Frobenius reciprocity and Mackey decomposition.  $\blacksquare$ 

Observe that the F-algebra  $\text{End}_E(V)$  is commutative if its dimension is at most 2. If  $\dim_F \text{End}_{\mathcal{E}}(V) = 1$ , V is absolutely indecomposable such that no proper submodule is an epimorphic image of V. In case  $V = F<sup>E</sup>$  is a

permutation module, i.e., U is the trivial  $FE<sub>0</sub>$ -module F, and G acts 2transitively,  $\text{End}_F(V)$  always is of F-dimension 2.

Write  $e_i = e_0 \otimes x_i$  for some right transversal  $(x_i)$  to  $E_0$  in E. If  $U^* = \langle e_0^* \rangle$ is the dual module of U,  $(U^*)^E$  can be viewed as the dual module  $V^*$ , with dual basis  $B^* = (e_0^* \otimes x_i)$ . The duality  $W \mapsto W^{\perp}$  from V to  $V^*$  preserves Einvariance. So every code  $(V, B, C)$  invariant under E corresponds to an Einvariant code  $(V^*, B^*, C^{\perp})$ . If  $U = U^*$  then we may identify  $(V, B)$  and  $(V^*, B^*)$ . This is familiar in case  $V = F^E$  is a permutation module. Clearly  $U = U^*$  whenever the corresponding character is of order at most 2 (e.g.,  $|F| \leqslant 3$ ).

(2.4) LEMMA. Suppose that  $\alpha$  is an automorphism of E normalizing  $E_{\alpha}$ . Let  $B_{\alpha} = (e_0 \otimes x_i^{\alpha^{-1}})$ . Then  $e_0 \otimes x_i \mapsto e_0 \otimes x_i^{\alpha^{-1}}$  defines a monomial isomorphism  $(U^E, B) \rightarrow ((U_{\alpha})^E, B_{\alpha})$  which gives a 1-1 correspondence between E-invariant codes.

*Proof.* It is immediate that  $B_a$  is a basis of  $\overline{V}=(U_a)^E$ . One checks that  $e_0 \otimes x_i \mapsto e_0 \otimes x_i^{\alpha^{-1}}$  defines a monomial isomorphism  $(V_\alpha, B) \to (\bar{V}, B_\alpha)$  of  $FE$ -modules. Moreover, the identity map  $(V, B) \rightarrow (V_a, B)$  is a monomial isomorphism respecting E-invariance of subspaces, by definition of  $V_{\alpha}$ .

Thus for instance the E-invariant codes in  $V = U^E$  and  $V^* = (U^*)^E$  are pairwise isomorphic if there is an automorphism  $\alpha$  of E normalizing  $E_{\alpha}$  and inverting the elements in  $E_0/C_{F_0}(U)$ .

 $(2.5)$  LEMMA. Let E be embedded in a finite group L. There exists a monomial action of L on  $V = U^E$  extending that of E if and only if there is a subgroup  $L_0$  with the following properties:

- (i)  $L = EL_0$  and  $E \cap L_0 = E_0$ ;
- (ii) U affords an  $FL_0$ -action extending that of  $E_0$ .

*Proof.* Straightforward.

If the field F of scalars is sufficiently large, condition (ii) in  $(2.5)$  is fulfilled exactly when there is a normal subgroup of  $L_0$ , with cyclic factor group, intersection  $E_0$  in the centralizer  $C_{E_0}(U)$ . We apply Lemma 2.5 mostly in the following situation: Suppose E is a subgroup of  $ML(C)$  for some indecomposable code C. If E induces a perfect permutation group  $G = G'$ , then E' is contained in any supplement L to  $F^*$  in  $ML(C)$ .

(2.6) LEMMA. Suppose char  $F = p > 0$  and  $\tilde{F}$  is a p-adic field with residue class field F. Then U can be lifted in a unique way to an  $FE_0$ -module  $\overline{U}$  preserving the order of the character, and  $V$  is the reduction of  $\tilde{\mathcal{V}} = \tilde{U} \otimes_{E_0} \tilde{F} E.$ 

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*Proof.* The order of the character  $\lambda$  afforded by U is prime to p. By Hensel's lemma there exists a unique  $\tilde{F}$ -character  $\tilde{\lambda}$  having the same order and lifting  $\lambda$ .

We finally give some comment concerning the fields of realization for codes. Let C be an FE-submodule of V and  $F_0$  a subfield of F. C can be written in  $F_0$  (with respect to E) if there is an E-invariant code  $(V_0, B_0, C_0)$ over  $F_0$  such that tensoring with F yields a code isomorphic to  $(V, B, C)$ .

(2.7) LEMMA. Suppose F is a splitting field for E of characteristic  $p > 0$ . Let  $\varepsilon$  be a root of unity in F such that all values of the (Frobenius) characters of the composition factors of  $V = U<sup>E</sup>$  and of the character afforded by U are powers of  $\varepsilon$ . If every semisimple section of V is multiplicity-free, then every code occurring as a submodule of  $V$  can be written in  $F_0 = \mathbb{F}_p(\varepsilon)$ .

*Proof.* Let  $V_0 = U_0 \otimes_{E_0} F_0 E$ , where  $U_0$  affords the  $F_0$ -character satisfying  $F \otimes U_0 = U$ . Since Schur indices over finite fields are 1, every composition factor of the  $F_0E$ -module  $V_0$  is absolutely irreducible. Using that the Jacobson radical  $J(FE) = F \otimes J(F_0E)$  we may conclude that  $W_0 \mapsto F \otimes W_0$ is an isomorphism from the lattice of submodules of  $V_0$  to that of  $V$ .

A corresponding result holds in characteristic 0 if the relevant Schur indices are 1, e.g., when  $G$  is 2-transitive (2.3).

## 3. PROJECTIVE PERMUTATION REPRESENTATIONS

In order to construct indecomposable codes we may use Schur's theory of projective representations. For the theoretical background we refer to [15].

(3.1) THEOREM. Let  $(V, B, C)$  an indecomposable code over F admitting a permutation group  $(G, \Omega)$ , where B is indexed by  $\Omega = \{0, ..., n-1\}$ . If  $Ext(G/G', F^*) = 0$ , every stem cover E of G affords an action on C inducing  $(G,\Omega).$ 

*Proof.* For each  $g \in G$  choose a preimage  $x_g$  in  $ML(C)$ . Then, by Theorem 1.2,  $g \mapsto x_g$  is a projective representation in the sense of Schur. Since  $\text{Ext}(G/G', F^*) = 0$  by hypothesis, there exists a homomorphism  $\varphi: E \to ML(C)$  making the diagram



commutative  $[15,$  Proposition V.5.5].

(3.2) COROLLARY. Assume  $(G, \Omega)$  is a primitive permutation group and F a field such that  $Ext(G/G', F^*) = 0$ . Let E be a stem cover of G and  $E_0$  be the inverse image in E of the stabilizer  $G_0$ . Induce up to E all 1-dimensional  $FE<sub>0</sub>$ -modules. Then the submodules of the resulting FE-modules provide for a complete list of codes over F admitting  $(G, \Omega)$  as permutation group.

*Proof.* Apply Proposition 2.1 and Theorems 1.3 and 3.1.

(3.3) Remark. Let  $\overline{F}$  be an algebraic closure of F. If  $(V, B, C)$  is a code over F, then tensoring with  $\overline{F}$  gives a code  $(\overline{V}, \overline{B}, \overline{C})$  over  $\overline{F}$ . If C admits  $(G, \Omega)$  then so does  $\overline{C}$ . Note that  $Ext(G/G', \overline{F^*}) = 0$  since  $\overline{F^*}$  is a divisible group. Therefore we may carry out the program of (3.2) over  $\bar{F}$  and then check whether the resulting codes can be written in  $F$  or not. Here Lemma 2.7 will be useful.

Clearly Ext $(G/G', F^*) = 0$  if  $G = G'$ . In case F is finite Ext $(G/G', F^*) = 0$ precisely when  $|G/G'|$  and  $|F^*|$  are relatively prime. Then every central Frattini extension  $A \rightarrow E \rightarrow G$  with A of exponent dividing  $|F^*|$  must be a stem extension, hence an epimorphic image over  $G$  of any stem cover of  $G$  $[15, Proposition V.5.5]$ . Therefore only that part of the Schur multiplier  $H_2(G) = H_2(G, Z)$  of G will be relevant in (3.1) and (3.2) which is of exponent dividing  $|F^*|$ .

The passage to an algebraic closure can be avoided sometimes, even when  $|G/G'|$  is not coprime to  $|F^*|$ :

 $(3.4)$  Proposition. Suppose  $(V, B, C)$  is an indecomposable code over the finite field F admitting a transitive permutation group  $(G, \Omega)$  of degree n. Assume the greatest common divisor of n,  $|G/G'|$ , and  $|F^*|$  is 1. Let A be the  $\pi$ -component of  $H_2(G)$ , where  $\pi$  is the set of primes dividing n. Then every stem extension  $A \rightarrowtail E \rightarrowtail G$  affords an action on C inducing (G,  $\Omega$ ).

*Proof.* Let L be the inverse image in  $ML(C)$  of G and  $L_0$  that of  $G_0$ . By Proposition 2.1 we may assume  $V = U^L$  where  $U = \langle e_0 \rangle$  is a 1-dimensional  $FL_0$ -module and  $B = (e_0 \otimes x_i)$  for some right transversal  $(x_i)$  to  $L_0$  in L. By Theorem 1.2, L is a central extension of G by  $F^*$ . It is immediate that  $L_0 = F^* \times C_{L_0}(U)$ .

Let B be the  $\pi$ -component of  $F^*$ . By Gaschütz's splitting theorem [6, Hauptsatz I. 17.4] there exists a supplement K to  $F^*$  in L intersecting  $F^*$  in B. Since  $|G/G'|$  is relatively prime to |B| by hypothesis, we have  $Ext(G/G', B) = 0$ . Therefore there exists a homomorphism  $\varphi: E \to K$  such that



commutes  $[15,$  Proposition V. 5.5]. This completes the proof. (Alternatively one could argue showing that any minimal supplement to  $B$  in  $K$  is a stem extension.)  $\blacksquare$ 

Investigating codes along the lines given in Corollary 3.2 the following proposition will be useful.

(3.5) LEMMA. Let  $A \rightarrow E \rightarrow G$  be a stem extension of the finite group G and  $E_0$  be the inverse image in E of some subgroup  $G_0$  of G. Let  $m = |G:G_0|$ . Then:

(a)  $A \cap E'_0$  is an epimorphic image of  $H_2(G_0)$  containing the m-th powers of the elements in A.

(b) If the corestriction map  $H_2(G_0) \to H_2(G)$  is surjective,  $A \rightarrow E_0 \rightarrow G_0$  is again a stem extension.

*Proof.* In view of the 5-term exact sequence  $\vert 15$ , Sect. II.3, the injection  $E_0 \rightarrow E$  yields the commutative diagram

$$
H_2(E_0) \longrightarrow H_2(G_0) \longrightarrow A \longrightarrow E_0/E'_0 \longrightarrow G_0/G'_0 \longrightarrow 1
$$
  
\n
$$
\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow
$$
  
\n
$$
H_2(E) \longrightarrow H_2(G) \longrightarrow A \longrightarrow E/E' \longrightarrow G/G' \longrightarrow 1
$$

having exact rows. By assumption the transgression  $H_2(G) \rightarrow A$   $(A = H_1(A))$ is epimorphic. (It is an isomorphism if and only if  $E$  is a stem cover of  $G$ .) This proves (b) and the first part of (a).

Consider the transfer from E to  $E_0/E_0'$ . Since  $A \subseteq E' \cap Z(E)$ , any  $x \in A$  is mapped onto  $E'_0 = x^m E'_0$ . Thus  $x^m \in A \cap E'_0$ .

We shall illustrate the program of  $(3.2)$  by discussing some highly transitive groups  $(G, \Omega)$ , namely the alternating groups and the Mathieu groups.

### 4. ALTERNATING GROUPS

We construct, up to isomorphism, all codes of length  $n \geq 4$  admitting the alternating group  $\mathfrak{A}_n$ . Of course, the repetition code  $(\langle \sum e_i \rangle)$  and its dual even admit the symmetric group  $\mathfrak{S}_n$ . It turns out that for  $n \geq 7$  no further code occurs. For  $n \le 6$  we obtain some other codes which, however, are well known. It is easily seen that all non-trivial codes admitting  $\mathfrak{A}_3$  are isomorphic to the repetition code or its dual.

 $(4.1)$   $\mathfrak{A}$ 

According to  $(3.3)$  we start with an algebraically closed field F of scalars. Up to group isomorphism,  $E = SL(2,3)$  is the unique stem cover of  $\mathfrak{A}_4$ . Following (3.2) we have to determine the nontrivial submodules of  $U^E$  where U is any 1-dimensional  $FE_0$ -module  $(E_0$  being the preimage of a point stabilizer).

(a) Let char  $F = 2$ . There are three nonisomorphic 1-dimensional  $FE_0$ . modules F, U, and U\*, which are restrictions of FE-modules F, W, and  $W^*$ , respectively. The obvious module isomorphisms  $U^E \cong W \otimes F^E$ ,  $(U^*)^E \cong W^* \otimes F^E$  are monomial (w.r.t. natural bases). So it is enough to study the permutation module  $V = F^E$  (=  $F^{\mathfrak{A}_4}$ ). V has two distinct submodules  $C_1$  and  $C_1^{\perp}$  which are interchanged by  $\mathfrak{S}_4$ . These are the (isomorphic) extended  $QR$ -codes over F.

From Lemma 2.7 it follows that the codes  $C_1$  and  $C_2^{\perp}$  can be written pecisely in those fields containing a primitive third root of unity.

We claim that  $PML(C_1) = \mathfrak{A}_4$ . At a first glance this seems to be obvious since the permutation group  $\mathfrak{S}_4$  interchanges  $C_1$  and  $C_1^{\perp}$ . But we have to exclude that there is a monomial action of a stem cover of  $\mathfrak{S}_4$  on V fixing  $C_1$ (Theorem 3.1). Since  $H_2(\mathfrak{S}_4) = Z_2$  and char  $F = 2$ , we actually are reduced to the permutation action,

(b) Assume char  $F = 3$ . There are just three irreducible  $FE$ -modules of dimensions 1, 2, 3, which can be realized over  $\mathbb{F}_3$ . The permutation module splits into the repetition code and its (irreducible) dual. Inducing up to  $E$  the unique nontrivial  $FE_0$ -module  $U$  gives an indecomposable  $FE$ -module  $V = U<sup>E</sup>$  whose unique proper submodule  $C_2$  has dimension 2. V affords a monomial action of  $GL(2, 3)$  extending that of  $E = SL(2, 3)$  by (2.5).  $C_2$  is invariant under  $GL(2, 3)$  as follows from Clifford theory. Hence  $PML(C_2) = \mathfrak{S}_4$ . Of course,  $C_2$  can be written in  $\mathbb{F}_3$  and then represents the  $(4, 2)$  Hamming code.

 $(c)$  It remains to consider the semisimple situation. We make use of the character table of  $SL(2,3)$  [4, Theorem 38.1]. Let U be the 1-dimensional  $FE_0$ -module affording the unique character of order 2,  $V = U^E$ . Then  $V = C_3 \oplus C_4$ , where  $C_3$  and  $C_4$  are irreducible (but nonisomorphic) FEmodules of dimension 2. By (2.5) we can extend the monomial action of  $E$ on V to  $GL(2,3)$  in two different ways. In both cases  $GL(2,3)$  interchanges  $C_3$  and  $C_4$ , which are the (isomorphic) extended QR-codes over F.  $C_3$  and  $C_4$  can be written just in such fields containing a primitive third root of unity, because of  $(2.7)$ .

All induced FE-modules of interest are of type  $W \otimes V$  or  $W \otimes F^E$  where W is a 1-dimensional FE-module. But the permutation module  $F<sup>E</sup>$  only yields the repetition code and its dual.

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We assert that  $PML(C_1) = \mathfrak{A}_4$ . Assuming the contrary there is a monomial action on V of the stem cover  $L = GL(2, 3)$  of  $\mathfrak{S}_4$  leaving  $C_3$ invariant. The restriction defines a monomial action of  $E = SL(2, 3)$  leaving  $C_3$  invariant. The preceding discussion shows that L cannot fix  $C_3$ .

 $(4.2)$   $\mathfrak{A}$ ,

As  $H_2(\mathfrak{A}_5) = Z_2$  [12], by (3.2) and (3.5a) we only have to investigate monomial actions of  $\mathfrak{A}_5$  (with permutation part  $\mathfrak{A}_5$ ).

The permutation module  $F^{\mathfrak{A}_{5}}$  yields the repetition code and its dual. If the field  $F$  does not contain a primitive third root of unity we do not get a proper monomial action.

Assume  $F$  contains a primitive third root of unity. Then the nontrivial 1dimensional  $F\mathfrak{A}_4$ -modules induce up to an  $F\mathfrak{A}_5$ -module V and its dual  $V^*$ . V is absolutely irreducible when char  $F \neq 2$ . (Note that char  $F \neq 3$ .) When char  $F = 2$ , V has a unique composition series  $0 \subset C_5 \subset C_5 \subset V$  where  $C_5$  is a (5, 2)  $QR$ -code and  $C_5$  is the expurgated  $QR$ -code.

Since  $\mathfrak{A}_5$  has an automorphism normalizing  $\mathfrak{A}_4$  and inverting  $\mathfrak{A}_4/\mathfrak{A}_4'$ , by Lemma 2.4 the dual module  $V^*$  gives codes isomorphic to  $C_5$  and  $C_5'$ . We claim that  $PML(C_5) = \mathfrak{A}_5$ . Otherwise, by passage to an algebraic closure, we have a monomial action of a stem cover L of  $\mathfrak{S}_5$  (Theorem 3.1). Lemma 2.5 and the remark following it lead to the desired contradiction.

Observe that  $C_5$ , written over  $\mathbb{F}_4$ , is the 1-perfect (5, 3) Hamming code. I

 $(4.3)$   $\mathfrak{A}_{6}$ 

In view of  $(4.2)$  just the repetition code and its dual will occur when char  $F \neq 2$  or F does not contain a primitive third root of unity. So assume  $\mathbb{F}_4 \subseteq F.$ 

It is known that  $H_2(\mathfrak{A}_6) = Z_6$  [12]. Since char  $F = 2$  we only have to investigate te monomial actions of the 3-fold cover  $A \rightarrow E \rightarrow \mathfrak{A}_6$  ( $|A| = 3$ ; E is the so-called Valentiner group). Let  $E_0$  be the inverse image in E of a point stabilizer  $\mathfrak{A}_5$ . Since  $H_2(\mathfrak{A}_5) = Z_2$  we have  $E_0 = E_0' \times A$ . Inducing up to E the nontrivial 1-dimensional  $FE_0$ -modules gives an FE-module V and its dual  $V^*$ .

There is a noninner involutory automorphism  $\bar{\alpha}$  of  $\mathfrak{A}_6$  normalizing  $\mathfrak{A}_5$ . By [15, Proposition V. 5.5]  $\bar{\alpha}$  can be lifted to an automorphism  $\alpha$  of E. Since  $\mathfrak{S}_6 = \langle \bar{a}, \mathfrak{A}_6 \rangle$  and  $H_2(\mathfrak{S}_6) = Z_2$ ,  $\alpha$  cannot centralize A. Therefore  $\alpha$  inverts  $E_0/E_0'$ . By Lemma 2.4 every E-invariant code in  $V^*$  is isomorphic to one in V. V contains a unique proper submodule  $C_6$ , the extended (6, 3) QR-code.

E has a 2-transitive subgroup  $H \cong PSL(2, 5)$  such that  $V = F<sup>H</sup>$  as an FHmodule.  $F^H$  has two distinct submodules  $C_6$  and  $C_6^{\perp}$  of dimension 3 being interchanged by  $PGL(2, 5)$  (see also Theorem 6.2 below). Since  $PGL(2, 5)$ supplements  $\mathfrak{A}_6$  in  $\mathfrak{S}_6$ , we may conclude from (2.5) that  $PML(C_6) = \mathfrak{A}_6$ .

(4.4) THEOREM. Let C be a nontrivial  $(n, k)$ -code over F. If  $n \ge 7$  and  $PML(C) \supseteq \mathfrak{A}_n$ , then C is isomorphic to the repetition code or its dual.

*Proof.* The permutation module for  $\mathfrak{A}_n$ , has the repetition code and its dual as unique proper submodules [7]. The theorem is established by showing that every monomial action of the stem cover E of  $\mathfrak{A}_n$  is the natural permutation action of  $\mathfrak{A}_n$ .

Let  $E_0$  be the inverse image in E of a stabilizer  $\mathfrak{A}_{n-1}$  (fixing the letter  $n-1$ ). We have to show that  $E_0 = E'_0$ . This will be a consequence of the fact that the corestriction  $H_2(\mathfrak{A}_{n-1}) \to H_2(\mathfrak{A}_n)$  is epimorphic, because of Lemma 3.5(b). In order to prove this we make use of Schur's work  $[12]$ . One knows that  $H_2(\mathfrak{A}_7) = Z_6$  and  $H_2(\mathfrak{A}_n) = Z_2$  for  $n \ge 8$ . However, we need more details and have to examine Schur's arguments more closely.

Consider the Moore presentation  $R \rightarrow L \rightarrow \mathfrak{A}_n$  of  $\mathfrak{A}_n$ , L being free on  $x_1,...,x_{n-2}$  and R generated as an L-group by  $x_1^3, x_i^2$  for  $2 \le i \le n-2$ ,  $(x_ix_{i+1})^3$  for  $1 \leq i \leq n-3$ , and  $(x_ix_j)^2$ , where  $1 \leq j < i-1$ ,  $i \leq n-2$ . The explicit presentation is given by  $x_1 \mapsto (0 \ 1 \ 2), x_i \mapsto (0 \ 1)(i, i+1)$  for  $2 \le i \le n-2$ . Setting  $T = \langle x_1, ..., x_{n-3} \rangle$  and  $S = T \cap R$  we obtain a free presentation  $S \rightarrow T \rightarrow \mathfrak{A}_{n-1}$ , and the corestriction  $H_2(\mathfrak{A}_{n-1}) \rightarrow H_2(\mathfrak{A}_n)$  is the natural map  $T' \cap S/[T, S] \rightarrow L' \cap R/[L, R].$ 

Schur [12, p. 117] proved that there is a word z in  $x_1, ..., x_4$  such that  $|z|L, R$  generates  $L' \cap R/[L, R]$ . (One may take, for instance,

$$
z = x_1 x_2 x_1 x_2 x_1^{-2} x_2^{-3} x_3 x_4^{-1} x_3^{-1} x_1^{-1} x_3^{-1} x_1^{-1} x_4 x_1 x_4.
$$

Since  $n \geq 7$  we have also  $T' \cap S/[T, S] = \langle z[T, S] \rangle$ . This gives the desired  $conclusion.$ 

 $(4.5)$  Remark. It readily follows from Lemma 3.5 that the corestriction  $H_2(\mathfrak{A}_{n-1}) \to H_2(\mathfrak{A}_n)$  is epimorphic if n is odd,  $n \geqslant 5$ . Using this information one can establish Theorem 4.4 inductively by shortening the codes of even length.

### 5. MATHIEU GROUPS

For a discussion of the Mathieu groups  $\mathfrak{M}_n$ , we refer to Conway [3] and Lüneburg  $[8]$ . It is known that the (extended) Golay codes admit Mathieu groups as permutation groups  $[3]$ . We will show that there are no further interesting codes with this property.

 $(5.1) \mathfrak{M}_{11}$ 

As  $H_2(\mathfrak{M}_{11}) = 0$  [2], by (3.2) we have to investigate induced modules  $V = U^{\mathfrak{M}_{11}}$ , where U is a 1-dimensional  $FW_{10}$ -module over some field F. The

commutator factor group of the point stabilizer  $\mathfrak{M}_{10}$  has order 2. The permutation module  $F^{\mathfrak{M}_{11}}$  yields the repetition code and its dual [7, Satz 4]. In case char  $F \neq 2$  we have a nontrivial 1-dimensional  $FW_{10}$ -module U. Let  $V= U^{\mathfrak{M}_{11}}$  for that U.

We first show that V is irreducible if char  $F \neq 3$ . Let  $F = \mathbb{Q}$ . Applying (2.3) we see that the  $\mathbb{Q} \mathfrak{M}_{11}$ -module V is absolutely irreducible. In view of  $(2.6)$  we have to verify that V remains irreducible modulo 11 and 5  $(|\mathfrak{M}_{11}| = 11 \cdot 10 \cdot 9 \cdot 8)$ . This is clear mod 11 since V belongs to an 11-block of defect 0 |4 Sect. 62]. Looking up the character table for  $\mathfrak{M}_{11}$  in [2] one realizes the character decomposition

$$
\chi_{44} = \chi_1 + \chi_{11} + \chi_{16} + \chi_{16}^*
$$

on 5-regular elements  $(\chi_{11} = \text{character of } V; \chi_n(1) = n)$ . There are five 5blocks of defect 0 and 9 conjugacy classes of 5-regular elements. Since the Sylow 5-subgroups of  $\mathfrak{M}_{11}$  have order 5, the decomposition numbers are 0 or 1 by Brauer-Dade [4, Theorem 68.1]. We may conclude that the restrictions to 5-regular elements of  $\chi_1$ ,  $\chi_{11}$ ,  $\chi_{16}$ ,  $\chi_{16}^*$  are just the irreducible Brauer characters in the principal 5-block for  $\mathfrak{M}_{11}$ . In particular, V is irreducible as well when  $F$  is a field of characteristic 5.

So let char  $F = 3$ . Then V is a uniserial  $FW_{11}$ -module with composition series  $0 \subset C_{11} \subset C_{11} \subset V$ , where  $C_{11}$  is the ternary Golay code of dimension 6. (Note that  $C_{11}^{\perp}$  is the expurgated Golay code.) All these facts can be established using the information given in Conway  $[3]$ .  $C_{11}$  is absolutely irreducible because its dimension is the prime 5 [4, Theorem 24.6]. From (2.3) it follows that  $C_{11}^{\perp}$  is not isomorphic to  $V/C_{11}$ which, in fact, is the dual module of  $C_{11}$ .

 $(5.2.)$   $\mathfrak{M}_1$ ,

By [2]  $H_2(\mathfrak{M}_{12}) = Z_2$ . Let E be the stem cover of  $\mathfrak{M}_{12}$  and  $E_0$  the inverse image in E of a point stabilizer  $\mathfrak{M}_{11}$ . Then  $E_0$  is a direct product of  $Z_2$  and a copy of  $\mathfrak{M}_{11}$ . As before the permutation module gives only the repetition code and its dual. If char  $F \neq 2$ , the unique nontrivial 1-dimensional  $FE_{0}^$ module induces up to an  $FW_{12}$ -module V. When char  $F = 3$ , V has a unique proper submodule  $C_{12} = C_{12}$ , the extended ternary Golay code of dimension 6.  $C_{12}$  is an absolutely irreducible  $FW_{12}$ -module being not isomorphic to its dual  $V/C_{12}$ . If char  $F \neq 3$ , V is irreducible by (5.1).

 $(5.3)$   $\mathfrak{M}_{22}$ 

It is known that  $H_2(\mathfrak{M}_{22}) = Z_{12}$  [18]. Let  $A \rightarrow E \rightarrow \mathfrak{M}_{22}$  be the stem cover of  $\mathfrak{M}_{22}$  and  $E_0$  be the inverse image in E of a point stabilizer  $\mathfrak{M}_{21} = PSL(3, 4)$ . By Lemma 3.5(a), 3 divides  $|A \cap E_0|$ . Assuming  $A \nsubseteq E_0$ . we get a character of order 2 of  $E_0/E_0'$ , producing a faithful complex module M of dimension 22 for the 2-fold proper covering  $\bar{E}$  of  $\mathfrak{M}_{22}$ . By (2.3) M has at most two irreducible components. But  $\overline{E}$  has only faithful irreducible complex representations of degree 10 and of degree larger than 55  $[2, p. 304]$ and  $\mathfrak{M}_{22}$  only those of degree 21 and at least 55 [2, p. 744]. This forces  $A \subseteq E_0 = E_0$ . (As a matter of fact, we see that 12 divides  $|H_2(\mathfrak{M}_{21})|$ ; it is known that  $H_2(\mathfrak{M}_{21}) = Z_{12} \times Z_4$ .)

In order to determine the codes admitting  $\mathfrak{M}_{22}$  we therefore have to discuss the permutation module  $V = F^{\mathfrak{M}_{22}}$ . If char  $F \neq 2$ , we just obtain the repetition code  $C_F$  and its dual  $C_F^{\perp}$  [7, Satz 4]. Suppose  $F = \mathbb{F}_2$ . It is well known that  $\mathfrak{M}_{22}$  leaves invariant a (22, 12)-code  $C_{22}$  over  $\mathbb{F}_2$  which is obtained by shortening the binary Golay code  $C_{23}$ . We have  $V = C_F^{\perp} + C_{22}$ and  $C_F^{\perp} \cap C_{22} = C_F \oplus C_{22}^{\perp}$ .  $\mathfrak{M}_{22}$  acts trivially on  $C_{22}/C_{22}^{\perp}$  and, as 11 does not divide  $2^a - 1$  for  $a < 10$ , irreducibly on  $C_{22}^{\perp} \cong (V/C_{22})^*$ . From [7, Satz 4] it follows that  $C_F^{\perp}/C_F$  is (absolutely) indecomposable.

Hence all interesting  $\mathfrak{M}_{22}$ -invariant codes (over  $\mathbb{F}_2$ ) are situated between  $C_{22}$  and  $C_{22}^{\perp}$ . We claim that  $C_{22}^{\perp}$  is an absolutely irreducible  $\mathbb{F}_2 \mathfrak{M}_{22}$ -module. Note first that  $C_{22}^{\perp}$  is the set of all vectors in  $C_{22}$  of weights 0, 8, 12, 16.  $C_{22}^{\perp}$ contains 77 vectors of weight 16 complementary to the blocks of the system  $S(3,6,22)$ . The stabilizer T in  $\mathfrak{M}_{22}$  of a vector of weight 16 is a maximal subgroup having two orbits of length 6 and 16 on the 22 letters [3, Table 3]. Hence T fixes only 2 vectors in  $C_{22}^{\perp}$ . It follows  $\text{End}_{\mathfrak{M}_{22}}(C_{22}^{\perp}) = \mathbb{F}_2$ .

Thus the situation is the same for  $F \supseteq F$ ,.

 $(5.4) \mathfrak{M}_{23}$ 

As  $H_2(\mathfrak{M}_{23})=0$  [2] and  $\mathfrak{M}_{22}=\mathfrak{M}'_{22}$ , we just have to investigate the permutation module  $V = F^{\mathfrak{M}_{23}}$ . As before only the case where char  $F = 2$  is interesting. Then  $V = C_F \oplus C_F^{\perp}$ , where  $C_F^{\perp}$  has a unique proper submodule  $C_{23}^{\perp}$ .  $C_{23} = C_F \oplus C_{23}^{\perp}$  is the "binary" (23, 12) Golay code. 

## $(5.5) \mathfrak{M}_{24}$

As  $H_2(\mathfrak{M}_{24}) = 0$  [2] and  $\mathfrak{M}_{23} = \mathfrak{M}'_{23}$ , again only the permutation module  $V = F^{\mathfrak{M}_{24}}$  is of interest, where char  $F = 2$ . Then V has a unique composition series  $0 \subset C_F \subset C_{24} = C_{24}^{\perp} \subset C_F^{\perp} \subset V$ . Here  $C_{24}$  is the extended "binary" (24, 12) Golay code. The indecomposability of  $C_F^{\perp}/C_F$  again follows from [7, Satz 4]. By weight consideration (over  $F_2$ ), using the fact that  $\mathfrak{M}_{24}$  acts transitively on the set of dodecads [3], one realizes that  $C_{24}$  is (absolutely) indecomposable.  $\blacksquare$ 

### (5.6) Permutation Groups

We determine the permutation groups  $PML(C_n)$ . Clearly the Golay codes are (extended)  $QR$ -codes. We will see in Section 6 that an extended  $QR$ -code of length  $p + 1$  does not admit PGL(2, p). Since  $\mathfrak{M}_n$  is a maximal subgroup of  $\mathfrak{A}_n$ , we may conclude from Theorem 4.4 that  $PML(C_n) = \mathfrak{M}_n$  for

 $n = 12, 24$ . This holds also for  $n = 11, 23$ , because  $\mathfrak{M}_{11}$  and  $\mathfrak{M}_{23}$  have no outer automorphisms. Of course, these facts are well known [3, Theorems 2.4 and  $3.6$ ].

The situation is different for  $\mathfrak{M}_{22}$ . Aut $(\mathfrak{M}_{22})$  is a permutation group of degree 22 [8, 12.5], having  $\mathfrak{M}_{22}$  as a normal subgroup of index 2 [3]. In fact, Aut( $\mathfrak{M}_{22}$ ) is induced by the normalizer  $N_{\mathfrak{M}_{24}}(\mathfrak{M}_{22})$  and so leaves invariant  $C_{22}$  and  $C_{22}^{\perp}$ . There is only one  $\mathfrak{M}_{22}$ -invariant code of dimension 11 which admits Aut $(\mathfrak{M}_{22})$ , namely,  $C_{22} \cap C_F^{\perp}$ . Since Aut $(\mathfrak{M}_{22})$  is a maximal subgroup of  $\mathfrak{S}_{22}$ , (4.4) gives  $PML(C_{22}) = Aut(\mathfrak{M}_{22})$ .

### 6. EXTENDED QR-CODES

By the Gleason-Prange theorem [1, Theorem 3.1] the extended  $QR$ -codes of length  $p + 1$ , p an odd prime, admit  $PSL(2, p)$ . We will characterize these codes by the property that their permutation group contains  $PSL(2, p)$  but not PGL(2, p). The case  $p = 3$  is already handled in (4.1) so that we may assume  $p \ge 5$ . Then  $PSL(2, p)$  is simple.

 $G = PSL(2, p)$  is 2-transitive on  $p + 1$  letters, the points of the projective line  $\Omega$  over  $\mathbb{F}_p$ . It is known that  $E = SL(2, p)$  is the unique stem cover of G [6, Satz  $V. 25.7$ ]. Because of Theorem 3.1 we have to investigate monomial actions of E. Write  $a = \begin{pmatrix} v & 0 \\ 0 & v^{-1} \end{pmatrix}$  for some generator v of  $\mathbb{F}_p^*$ , and let  $u = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$ ,  $c = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ . We have  $u^{-1}au = a^{-1}$ .  $E_0 = \langle a, c \rangle$  is the normalizer of the Sylow p-subgroup  $S = \langle c \rangle$  of E.  $H = \langle a \rangle$  complements S in  $E_0$ .  $E_0$  is the inverse image of a point stabilizer  $G_0$  (fixing  $\infty$ ). The normalizer  $N = N_E(H)$ is generated by a and u. For any  $x \in H$  with  $x^2 \neq 1$ ,  $C_F(x) = H$  and  $N_E(\langle x \rangle) = N.$ 

The above notation is fixed through (6.1), (6.2).

(6.1) LEMMA. Suppose U is a 1-dimensional  $FE_0$ -module affording a character  $\lambda$  of order greater than 2; let  $V = U^E$ . Then  $\text{End}_E(V)$  is of Fdimension 1. If char  $F \neq p$ , V is absolutely irreducible.

*Proof.* Clearly  $u \in E - E_0$  and  $H = E_0 \cap E_0^{u^{-1}}$ . Since u inverts the elements of H and  $\lambda$  has order greater than 2, by (2.3) End<sub>r</sub>(V) has dimension 1. Hence V is absolutely irreducible if char  $F$  does not divide  $|E| = (p + 1)p(p - 1)$ . We may assume that F is an algebraically closed field such that char  $F = q$  divides  $p^2 - 1$ . Note that the order of  $\lambda$  is prime to q.

Let  $\tilde{F}$  be a q-adic field with residue class field F. According to Lemma 2.6 we can lift V in a natural way to an  $\overline{FE}$ -module  $\overline{V}$ .  $\overline{V}$  is absolutely irreducible. If q is odd and a divisor of  $p + 1$ , the order of a Sylow qsubgroup of E divides  $p + 1$ . If  $q = 2$ ,  $G = PSL(2, p)$  operates on  $\tilde{V}$ , and  $p + 1$  is divisible by the order of a Sylow 2-subgroup of G when  $p \equiv 3$ (mod 4). In these cases V and  $\tilde{V}$  belong to a q-block of defect 0 and thus V is irreducible  $[4, Sect. 62].$ 

Assume therefore that either  $q = 2$  and  $p \equiv 1 \pmod{4}$  or q is odd and a divisor of  $p-1$ . Let  $D = \langle x \rangle$  be a Sylow q-subgroup of H. Denote by  $\chi$  the (ordinary) character afforded by  $\tilde{V}$ . From the character table of  $SL(2, p)$  [4, Theorem  $38.1$ ] one sees that

$$
\chi(xy) \equiv \lambda(y) + \lambda^*(y) \qquad \text{(mod } q\text{)}
$$

for all  $y \in H$ ;  $\lambda^*$  is the dual character to  $\lambda$ . Now  $H = C_R(x)$  since  $|D| > 2$ . If  $\lambda$ |H belongs to the block b of FH, then  $\chi$  belongs to the q-block  $B = b^E$  by Brauer's second main theorem  $[4,$  Theorem 63.2]. Clearly D is a defect group of b. As  $\lambda$  has order  $\neq 1,2$  and u inverts the elements of H, we have  $N_N(b) = H = C_F(D)$ . Applying Brauer's first main theorem  $\left[4\right]$ Theorems 64.10 and 58.3] shows that  $D$  is also a defect group of  $B$ . From [4, Theorem 68.1] it follows that V is the unique irreducible  $FE$ -module in the block B. This completes the proof.  $\blacksquare$ 

 $(6.2)$  THEOREM. Besides the repetition code and its dual, there are  $(up_to$ isomorphism) precisely the following proper codes admitting  $G = PSL(2, p)$ :

(i) If char  $F = 2$  and  $\mathbb{F}_4 \subseteq F$ , there is a  $(p+1, (p+1)/2)$ -code over F; it can be written in  $\mathbb{F}$ , if and only if  $p \equiv \pm 1 \pmod{8}$ .

(ii) If char F is different from 2 and p and  $(-1)^{(p-1)/2}$  p is a square in  $F^*$ , there is a (p + 1, (p + 1)/2)-code over F.

(iii) In case char  $F = p$  there exist  $(p + 1, k)$ -codes over F, one for each k between 2 and  $p-1$ .

The codes in (iii) admit  $PGL(2, p)$ ; the codes appearing in (i), (ii) are the extended QR-codes which do not admit  $PGL(2, p)$ .

*Proof.* Recall that  $PGL(2, p)$  is sharply 3-transitive on  $\Omega$ . It is immediate that  $GL(2, p)$  is a stem cover of  $PGL(2, p)$ . Every monomial action of  $E = SL(2, p)$  can be extended, in various ways, to  $GL(2, p)$ .

(i) char  $F = 2$ .

In view of  $(6.1)$  we only have to study the permutation module  $V = F<sup>G</sup>(= F<sup>E</sup>)$ . Assume first that F is algebraically closed. Let  $1 + \psi$  be the complex permutation character of  $G = PSL(2, p)$ . From the character table [4, Theorem 38.1] one obtains that there are irreducible characters  $\eta_1$ ,  $\eta_2$  (of G) of degree  $(p-1)/2$  such that

$$
\psi = 1 + \eta_1 + \eta_2
$$

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on 2-regular elements;  $\eta_1$  and  $\eta_2$  differ on 2-regular elements, and they have there values in  $\mathbb{Q}(\varepsilon)$ , where  $\varepsilon^2 = (-1)^{(p-1)/2} p$ .

Now  $G_0$  is a Frobenius group with kernel the image  $\overline{S}$  of S in G. Application of  $[6, Satz V.16.13]$  and Mackey decomposition shows that the restriction of  $\psi$  to  $G_0$  can be written as

$$
\psi|G_0 = 1 + \eta'_1 + \eta'_2,
$$

where  $\eta'_1$  and  $\eta'_2$  are different irreducible characters of degree  $|G_0/\overline{S}| =$  $(p-1)/2$ . Since  $\eta'_1, \eta'_2$  are in a 2-block of defect 0 for  $G_0$ , we may deduce that  $\eta_1$  and  $\eta_2$  remain irreducible as Brauer characters. Hence V has (unique) submodules  $C_F \subset C_F^{\perp}$  of dimensions 1 resp. p and  $M = C_F^{\perp}/C_F$  has two nonisomorphic composition factors of dimension  $(p-1)/2$ .

Viewing V as the permutation module for  $PGL(2, p)$ , M is irreducible. This follows, for instance, from the fact that the permutation character of a point stabilizer is of type  $1 + \gamma$ , where  $\gamma$  is irreducible of degree  $p - 1$  and so is in a 2-block of defect 0. Consequently (Clifford)  $M = \overline{C}_1 \oplus \overline{C}_2$  for some nonisomorphic irreducible FG-modules  $\overline{C}_i = C_i/C_F$ . As codes the C<sub>i</sub> are isomorphic since they are interchanged by  $PGL(2, p)$ . Clearly  $C_1$  and  $C_2$  are the (isomorphic) extended  $QR$ -codes over  $F$ . (For an alternate approach see  $[16]$  or  $[11]$ .)

Since  $\eta_1$ ,  $\eta_2$  have their values in the quadratic field  $\mathbb{Q}(\varepsilon)$ , from Lemma 2.7 it follows that the codes  $C_i$  can be written in the field  $\mathbb{F}_4$  and, by elementary properties of 2-adic squares, in  $F_2$  precisely when  $p \equiv \pm 1 \pmod{8}$ .

Finally, since char  $F = 2$ , the permutation action of G cannot be extended to a proper monomial action of  $GL(2, p)$ . Hence  $PGL(2, p) \not\subseteq PML(C_i)$ .

(ii) char  $F \neq 2$ , p.

The permutation module  $V = F^E$  now yields only the repetition code and its dual. This can be checked by applying [ 13, Corollary 21. (One can verify this also by means of [7, Satz 8] in case char F does not divide  $p + 1$ , and by a block theoretic argument otherwise.)

In view of (6.1) it remains to consider  $V = U^E$ , where U is the 1dimensional  $FE_0$ -module affording the unique linear character  $\lambda$  of order 2. Assume first that  $F$  is algebraically closed. In the semisimple situation from [4, Theorem 38.1] (and its proof) it follows that  $V = C_1 \oplus C_2$  for some irreducible FE-modules  $C_i$  of dimension  $(p + 1)/2$ . So let char  $F = q$  be an odd divisor of  $p^2-1$ .

There is a q-adic field  $\tilde{F}$ , with residue class field  $F$ , which is a splitting field for E. Lift V to an FE-module  $\tilde{V}$  as in (2.6). We already know that there are irreducible characters  $\xi_1, \xi_2$  of degree  $(p + 1)/2$  such that  $\xi_1 + \xi_2$  is the character of  $\tilde{V}$ . From the character table we infer that  $\xi_1$  and  $\xi_2$  differ on q-regular elements and have their values in  $\mathbb{Q}(\varepsilon)$ ,  $\varepsilon$  as in (i). We claim that  $\xi_1, \xi_2$  are irreducible also as Brauer characters mod q.

Let D be a Sylow q-subgroup of E. If q is a divisor of  $p + 1$ , |D| divides the degree of the  $\xi_i$  and we are done. Suppose next that q is a divisor of  $p-1$ . We may assume  $D \subseteq H$ . From [4, Theorem 38.1] once more we obtain that there is an irreducible character  $\chi$  of E (with values in  $\tilde{F}$ ), induced from a linear character of  $E_0$  of order  $2 |D|$ , such that

$$
\chi = \xi_1 + \xi_2
$$

on q – regular elements. Let B denote the q-block containing  $\chi$ , hence also  $\xi_1$ and  $\xi_2$ . Since  $H \supseteq D$  is cyclic and B is not of defect 0, application of [4, Theorem 68.1 shows that  $\xi_1$  and  $\xi_2$ , restricted to q-regular elements, are the unique irreducible Brauer characters in  $B$ .

Consequently  $V$  has two nonisomorphic composition factors of dimension  $(p+1)/2$ . By (2.3) dim<sub>F</sub> End<sub>F</sub>(V) = 2. Thus, as before,  $V = C_1 \oplus C_2$  for some irreducible FE-modules  $C_i$  of dimension  $(p + 1)/2$ .

In any case,  $C_1$  and  $C_2$  represent the extended QR-code over F. In fact,  $GL(2, p)$  interchanges  $C_1$  and  $C_2$  in any monomial action extending that of  $E$ . This also can be seen from the character table. From  $(3.1)$  and the remark following (2.5) we may conclude that the codes  $C_i$  do not admit  $GL(2, p)$ .

The codes  $C_i$  can be written in  $\mathbb{F}_q(\varepsilon)$  resp.  $\mathbb{Q}(\varepsilon)$ , where  $\varepsilon^2 = (-1)^{(p-1)/2} p$ . This follows from (2.7); (2.3) guarantees that the Schur index of  $\xi_i$  over  $\mathbb Q$ is 1.

(iii) char  $F = p$ .

By Brauer-Nesbitt there is, up to isomorphism, exactly one (absolutely) irreducible FE-module  $W_k$  of dimension k  $(1 \le k \le p)$  [6, V.5.13]. From [4, Theorem 71.3] one obtains that the various 1-dimensional  $FE_0$ -modules induce up to FE-modules  $V_k$  having a submodule  $C_k \cong W_k$  such that  $V_k/C_k \cong W_{p+1-k}$   $(1 \le k \le p-1)$ . Furthermore  $V_{(p+1)/2} = U^E$ , where U affords the character of  $E_0$  of order 2. Obviously  $V_1$  is the permutation module, and  $V_1 = C_1 \oplus C_1^{\perp}$ .

By (2.3), (6.1) End<sub>E</sub> $(V_k)$  has F-dimension 2 precisely when  $k = 1$  or  $k = (p + 1)/2$ . Since both composition factors of  $V_{(p+1)/2}$  are isomorphic,  $V_{(p+1)/2}$  must be indecomposable. This is immediate in the other cases. Hence  $C_k$  is the unique proper submodule of  $V_k$ ,  $k = 2,..., p - 1$ . The code  $C_k$  admits  $PGL(2, p)$  by Clifford theory.

The codes in (iii) are extensions of the optimal codes in characteristic  $p$ described by Assmus and Mattson [1, Sect. 2].

## $(6.3)$  Permutation Groups

Suppose C is an extended QR-code of length  $p+1$  over F and  $G = PML(C)$ . Then G contains  $PSL(2, p)$  but not  $PGL(2, p)$ . Only four cases are known where  $G \neq PSL(2, p)$ , namely,

(1)  $p=5$ ,  $F \supseteq F_4$ :  $G=\mathfrak{A}_6$ , (2)  $p = 7$ , char  $F = 2$ :  $G = Aff(3, 2)$ (3)  $p = 11$ , char  $F = 3$ :  $G = \mathfrak{M}_{12}$ , (4)  $p = 23$ , char  $F = 2$ :  $G = \mathfrak{M}_{24}$ .

It is conjectured that  $G = PSL(2, p)$  provided  $p > 23$ . We cannot settle this in generality, but here is some further evidence for its truth.

 $(6.4)$  THEOREM. Let C be an extended QR-code over F of length  $p + 1 \geq 8$ , and let G be a subgroup of PML(C) containing PSL(2, p). Then

- (i) G is a proper subgroup of  $\mathfrak{A}_{n+1}$ .
- (ii) If  $p > 7$  and  $G \neq PSL(2, p)$ , then G is 4-transitive and simple.

*Proof.* Let N be the normalizer in  $\mathfrak{S}_{n+1}$  of a Sylow p-subgroup S of PSL(2, p). N has order  $p(p-1)$  and contains a  $(p-1)$ -cycle. It follows that N supplements  $PSL(2, p)$  in  $PGL(2, p)$ , and  $\mathfrak{A}_{p+1}$  in  $\mathfrak{S}_{p+1}$ . Since G contains  $PSL(2, p)$  but not  $PGL(2, p)$ , the normalizer  $\overline{N} = N_G(S)$  is a subgroup of  $PSL(2, p)$ .

S is a Sylow p-subgroup of  $G \cap \mathfrak{A}_{p+1}$ . Hence from  $\overline{N} \subseteq PSL(2, p) \subseteq$  $G \cap \mathfrak{A}_{p+1}$  it follows  $G \subseteq \mathfrak{A}_{p+1}$  by the Frattini argument. Because of (4.4) G is a proper subgroup of  $\mathfrak{A}_{n+1}$ .

Now assume  $G \neq PSL(2, p)$  and  $p > 7$ . By Neumann [10, Theorem 2.1] then G is 4-transitive. Suppose  $M \neq 1$  is a normal subgroup of G. M cannot be regular  $[17,$  Theorem 11.3], hence is at least 3-transitive. This implies that  $S \subseteq M$  and  $PSL(2, p) \subseteq M$ , by simplicity of PSL(2, p). Moreover we have  $G = M\overline{N}$ , again by the Frattini argument. Now from  $\overline{N} \subseteq PSL(2, p)$  it follows  $G = M$ , as desired.  $\blacksquare$ 

A group G as in Theorem 6.4(ii) would be an "unknown" simple group, provided  $p > 23$ . Theorem 6.4(i) answers a conjecture of Rasala to the affirmative  $[11, p. 470]$ . It should be possible to establish this by more elementary arguments than those used in (4.4). (But the argumentation by Shaughnessy  $[13, p. 402]$  cannot work, as follows from  $[1,$  Theorem 2.21.) Under additional assumptions, Theorem 6.4 can be improved so that  $PML(C) = PSL(2, p)$ . For instance, this holds if  $p-2$  is a prime [10, Corollary 2.2, or if  $(p-1)/2$  is a prime and  $23 < p \leq 4079$ . (The latter result has been already stated in  $[1, p, 146]$ . But, as Rasala  $[11]$  noted, it depends on the validity of his conjecture, i.e., on Theorem 6.4(i).)

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